

1

Session 2: Types

COMP2221: Functional programming

Laura Morgenstern*

*laura.morgenstern@durham.ac.uk

COMP2221-Session 2: Types

Recap

- What are some differences between functional and imperative programming?
- Which programming model more closely mirrors the way computers execute?
- What are interpreters and compilers? (in very broad terms)
- What are side effects?
- Why can side effects easily introduce bugs?
- What is Haskell's syntax for identifier definition, function application and composition?
- Which list operations does the Prelude provide?

Definition (Type)

A type is a *collection* of related values.

A type can be described by specifying

- the set of data elements it covers and
- the operations it supports.

Example

- Bool the two logical values True and False.
- Integer -> Integer the set of all functions that take an Integer as input and produce an Integer as output.

Example in C/Java int a = 4; double a = 4; int b = 3; double b = 3; double c = a/b; double c = a/b;

Example in C/Java

int $a = 4;$	<pre>double a = 4;</pre>
--------------	--------------------------

int b = 3; double b = 3;

double c = a/b; double c = a/b;

Result depends on input types.

Since computers represent *everything* as sequences of bits, types are also required to interpret these bit patterns.

Motivation: Why do we need types?

- Mathematics and programming rely on the notion of types
- Tell us *how* to interpret a variable
- Provide restrictions on valid operations
- Are required to know what a bit pattern means

Implementation

Find the correct implementation of an operator.

Correctness

Check whether an operation on some data is valid and/or well-defined.

Check whether a code fragment is correct (type safety).

Documentation

Document the code's semantics for the reader.

COMP2221—Session 2: Types

- Strongly vs. weakly typed languages
- Statically vs. dynamically typed languages

Translators must check for type correctness

Definition (Statically typed language) Type safety is checked at translation time.

 \Rightarrow invalid types result in translation error

```
-- Invalid
foo :: a -> Int
foo f = 1 + f
```

Definition (Dynamically typed language)

Type safety is checked at run time.

 \Rightarrow invalid types only detected as soon as used

```
# Fine as long as f supports addition with a number
def foo(f):
    return 1 + f
```

• How does the translator determine the type of an expression?

Explicit annotation

Programmer annotates all variables with type information (e.g. C/Java)

Type inference

Translator *infers* the types of variables based on the operations used (e.g. Haskell)

Duck typing

Translator/runtime just tries the operation, if it succeeds, that was a valid type! (Python)

Haskell is

Strongly, statically typed.

 \Rightarrow every well-formed expression has exactly one type, these types are known at $\mathit{compile time}$

Attaching types

```
Haskell's notation for "e is of type T" is spelt
```

```
e :: T
-- False is of type Bool
False :: Bool
-- not is of type Bool -> Bool
not :: Bool -> Bool
```

What type does X have?

Every valid expression in Haskell must have a valid type.

You can ask GHCi what the type of an expression is with the command :type expr

```
Prelude> :type sum
sum :: Num a => [a] -> a
```

Demo time

Let's look at some types

• Functions have types in all programming languages, Haskell makes this particularly explicit

Functions of one argument "unary"

Map from one type to another

```
not :: Bool -> Bool
and :: [Bool] -> Bool
```

Functions of two arguments "binary"

Map from two types to another

```
add :: (Int, Int) -> Int
```

"add eats two Ints and returns an Int"

• Content

- Defined and motivated types
- Different concepts of typing (dynamic/static)
- Considered basic Haskell types
- Looked at list and tuple types
- Determined type of expressions with GHCi
- Considered types of functions
- Self-study
 - Tackle exercise "Types and Lists"